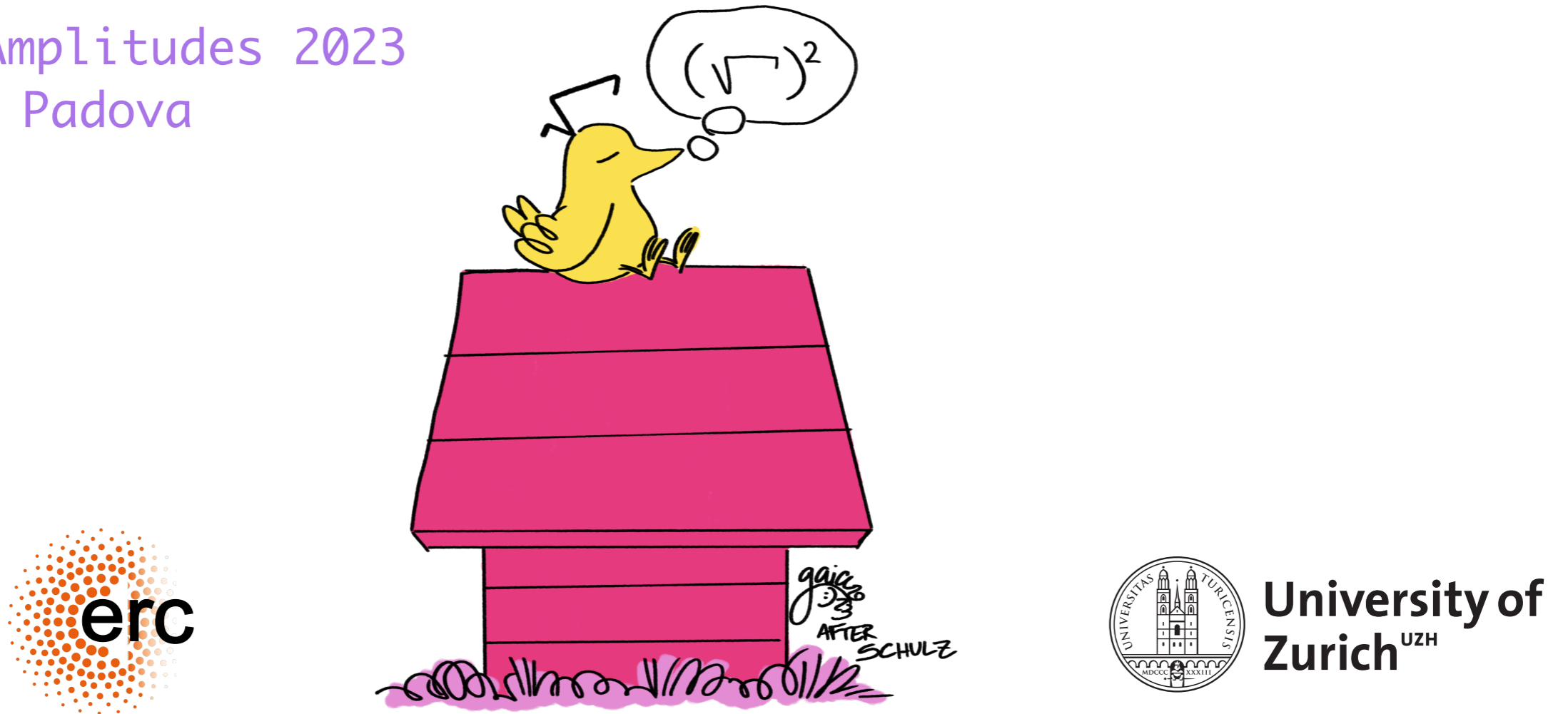


Intersection numbers & rational algorithms

MathemAmplitudes 2023
Padova



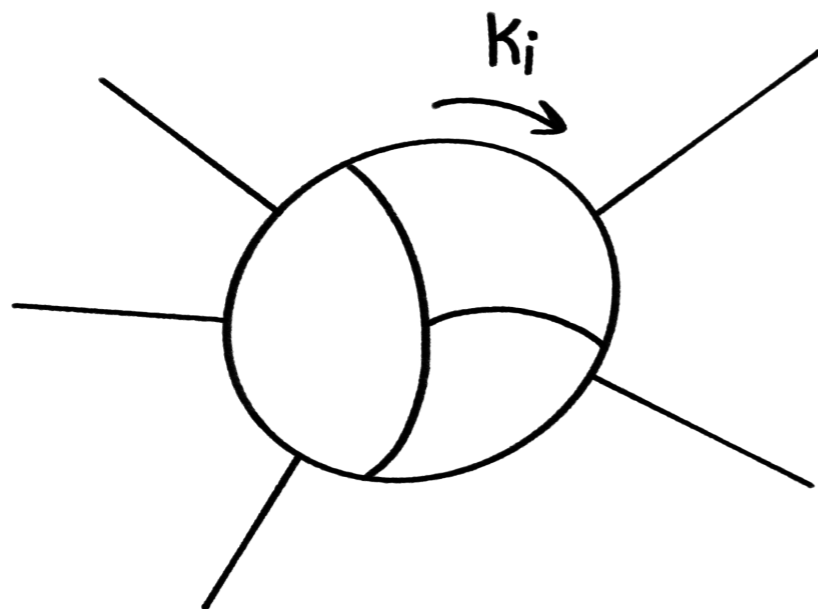
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Based on [arXiv:2304.14336](https://arxiv.org/abs/2304.14336)
in collaboration with Tiziano Peraro

Loop integrals

- * LEGO® blocks of perturbative QFT beyond tree level
- * Key ingredient of phenomenological predictions
- * Rich and interesting mathematical structures

$$I_i[\alpha_1, \dots, \alpha_n] = \int_{-\infty}^{+\infty} \left(\prod_{i=1}^{\ell} d^d k_i \right) \frac{1}{z_1^{\alpha_1} \dots z_n^{\alpha_n}}$$



inverse propagators
 $z_i = k_i^2 - m_i^2 + \text{ISPs}$



Not all are linearly
independent!

Reduction to master integrals

why?

- * Extremely large number of integrals contributing to an amplitude
- * Properties/symmetries of an amplitude manifest only after the reduction
- * Important for the calculation of the integrals

Reduction into a basis of linearly independent master integrals $\{G_j\} \subset \{I_j\}$

$$I_j = \sum c_{jk} G_k$$

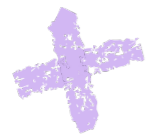
$\{G_j\}$ = minimal linearly independent set

Laporta algorithm

Feynman integrals in dimensional regularization obey linear relations, e.g. **Integration By Parts** identities

Chetyrkin, Tkachov (1981), Laporta (2000)

$$\int \left(\prod_{i=1}^{\ell} d^d k_i \right) \frac{\partial}{\partial k_i^\mu} \left(\frac{v_j^\mu}{z_1^{\alpha_1} \dots z_n^{\alpha_n}} \right) = 0, \quad v^\mu = \begin{cases} p_i^\mu = \text{external} \\ k_i^\mu = \text{loop} \end{cases}$$

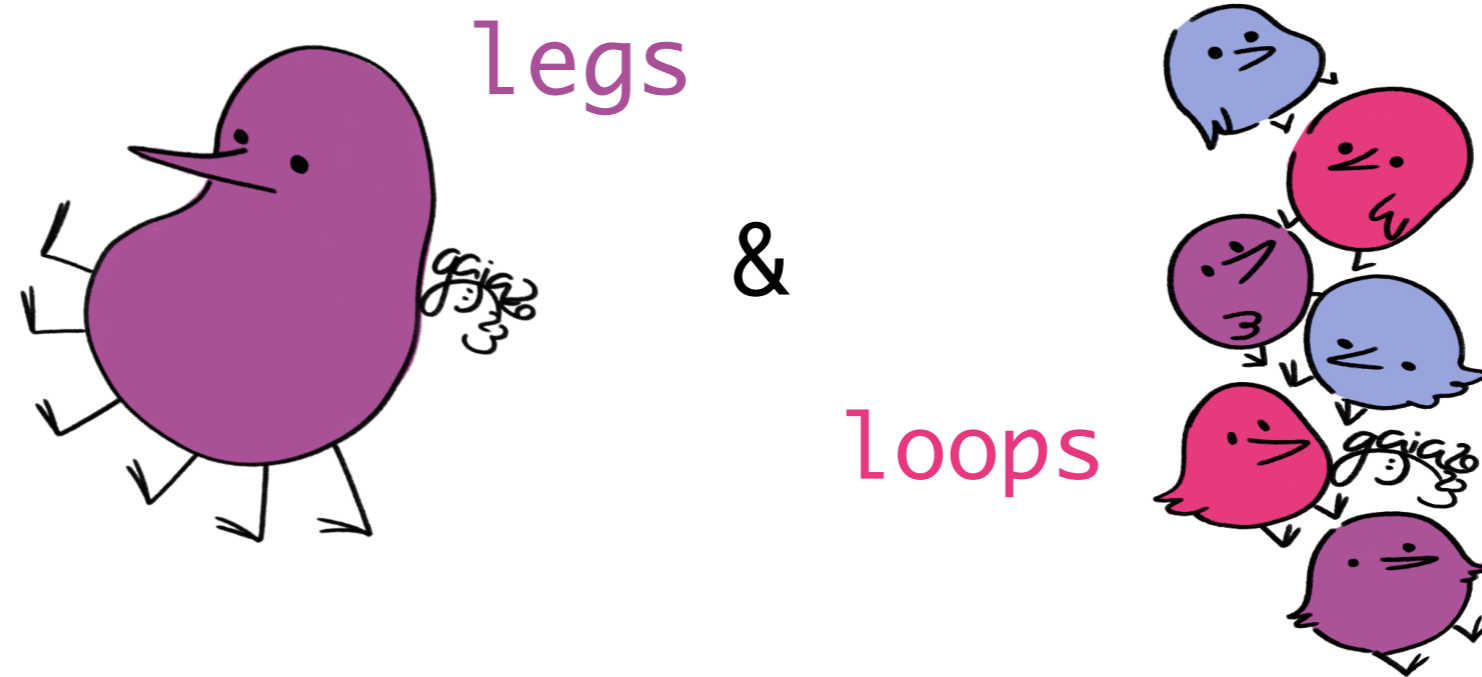


Lorentz Invariance ids, symmetry relations, ...

reduction as solution of a **large**
and **sparse** system of identities

Algebraic complexity

Processes with many



give rise to **HUGE** intermediate expressions

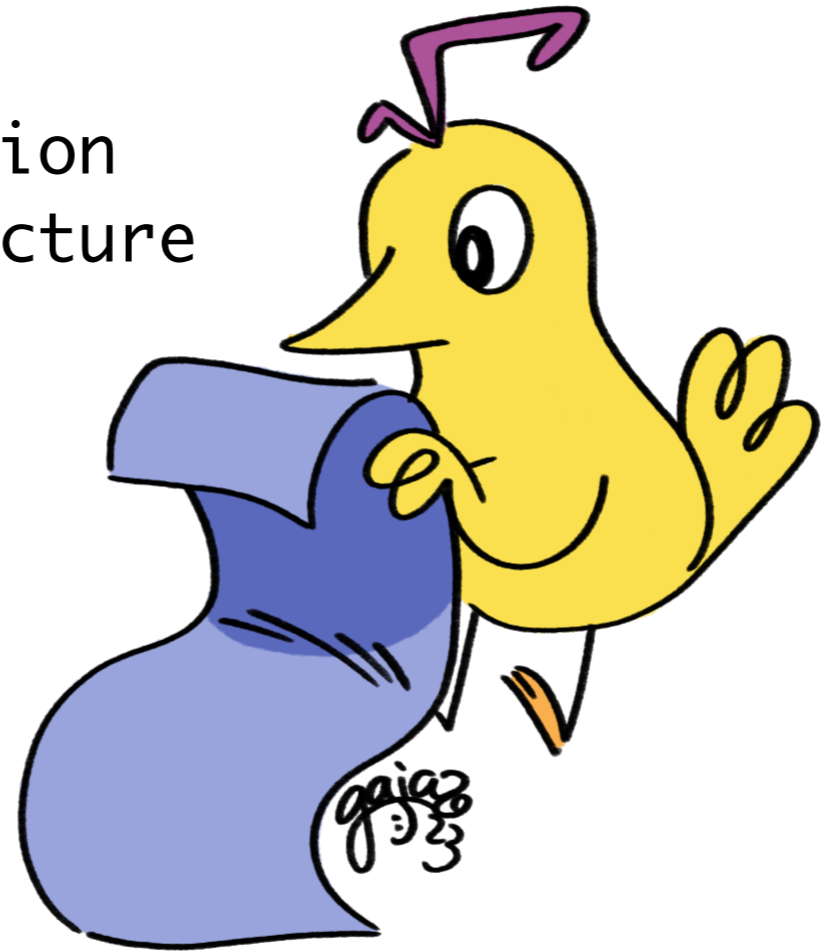
Drawbacks of Laporta procedure

- * Very large system \rightarrow computational bottleneck
- * Algebraic structure of FI not manifest

Looking for other ways...

Wishlist:

- * Allows for a direct decomposition
- * Exploits the vector space structure obeyed by Feynman integrals



One option:

Intersection theory

Framework

Vector space of n -folds integrals in $\mathbf{z} = (z_1, \dots, z_n)$

“right” integrals

$$|\varphi_R\rangle = \int dz_1 \dots dz_n \frac{1}{u(\mathbf{z})} \varphi_R(\mathbf{z})$$

“dual” or “left” integrals

$$\langle \varphi_L | = \int dz_1 \dots dz_n u(\mathbf{z}) \varphi_L(\mathbf{z})$$

with

• φ_L/φ_R rational functions • $u(\mathbf{z}) = \prod_j B(\mathbf{z})^{\gamma_j}$, $\begin{cases} \gamma_j \text{ generic exponents} \\ B_j \text{ polynomials} \end{cases}$

Intersection numbers:

scalar products between $\langle \varphi_L | \varphi_R \rangle$ Mastrolia, Mizera (2018)
left and right integrals

Vector space:

- * Dimension ν
- * Basis $|e_i^{(R)}\rangle$ and dual basis $\langle e_i^{(L)}|$
- * Scalar product: intersection number

Change of representation

Baikov change of vars

$$k_j \rightarrow z_j$$

Baikov (1996)

$$I[\alpha_1, \dots, \alpha_n] = \int \left(\prod_{i=1}^{\ell} d^d k_i \right) \prod_{j=1}^n \frac{1}{z_j^{\alpha_j}} \rightarrow \int dz_1 \dots dz_n B^\gamma \prod_{j=1}^n \frac{1}{z_j^{\alpha_j}} = |\varphi_R\rangle$$

$$|\varphi_R\rangle = \int dz_1 \dots dz_n \frac{1}{u(\mathbf{z})} \varphi_R(\mathbf{z})$$

with

$$u(\mathbf{z}) = B^{-\gamma} \prod z_j^{\rho_j}, \quad \gamma = \frac{d - E - L - 1}{2}$$

$$\varphi_R(\mathbf{z}) = \frac{1}{z_1^{\alpha_1} \dots z_n^{\alpha_n}}$$

analytic regulators
of $z_j \rightarrow 0$ singularities

Identifications

● $|\varphi\rangle$ generic vector

→ Feynman integral to reduce

● $\{|e_i^{(R)}\rangle\}_{i=1}^{\nu}$ basis vectors

→ master integrals

decomposition of integrals as

$$|\varphi_R\rangle = \sum_{i=1}^{\nu} c_i^{(R)} |e_i^{(R)}\rangle \quad c_i^{(R)} = \sum_{j=1}^{\nu} (\mathbf{C}^{-1})_{ij} \langle e_j^{(L)} | \varphi_R \rangle$$

where we introduce the metric

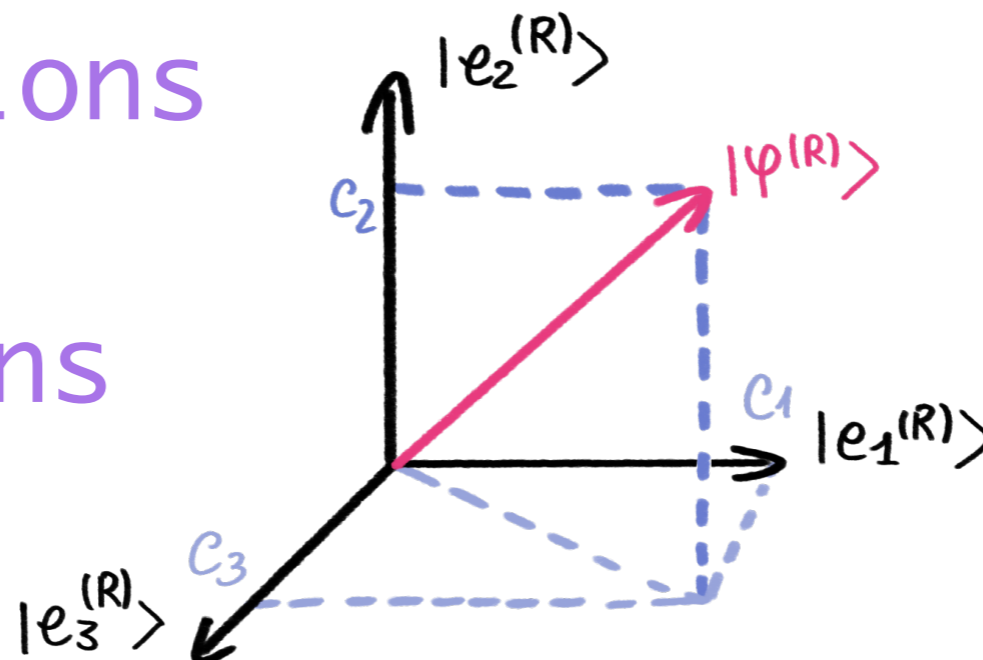
$$\mathbf{C}_{ij} = \langle e_i^{(L)} | e_j^{(R)} \rangle$$

● similar formulae for dual integrals

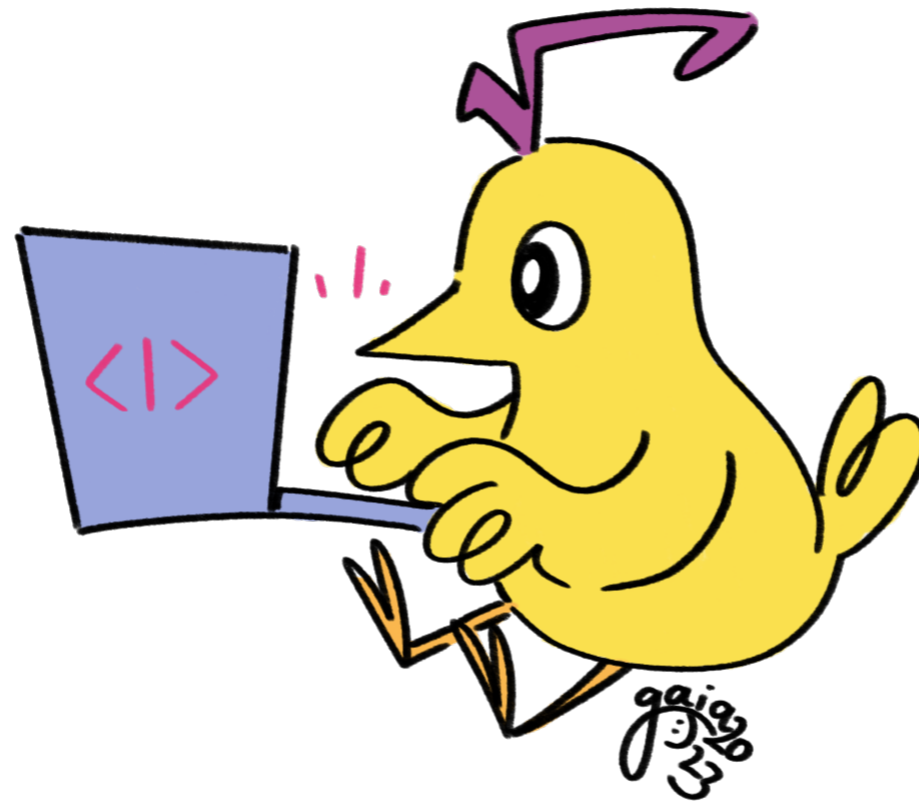
Decompositions

=

projections



Computation of intersection numbers



Univariate algorithm

Frellesvig et al. (2019)

We have 1-fold integrals
in the variable z

$$|\varphi_R\rangle = \int dz \frac{1}{u(z)} \varphi_R(z)$$

- univariate intersection numbers

$$\langle \varphi_L | \varphi_R \rangle = \sum_{p \in \mathcal{P}_\omega} \text{Res}_{z=p}(\psi \varphi_R)$$

- ψ is the local solution of

$$(\partial_z + \omega)\psi = \varphi_L, \quad \omega \equiv \frac{\partial_z u}{u}$$

- around each $p \in \mathcal{P}_\omega$

$$\mathcal{P}_\omega = \{z \mid z \text{ is a pole of } \omega\} \cup \{\infty\}$$

- ansatz around p

$$\psi = \sum_{i=\min}^{\max} c_i (z-p)^i + O\left((z-p)^{\max+1}\right)$$

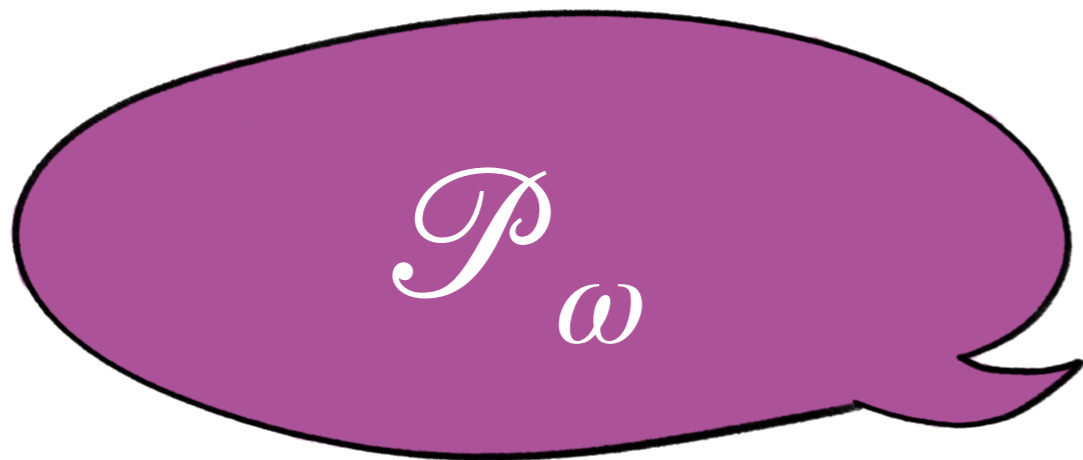
- * plug in the DE
- * solve for the c_i

$$\langle \varphi_L | \varphi_R \rangle = \text{rational}$$

BUT

- * Non-rational contributions in intermediate stages
- * Cancellations after sum over all residues

non-rational terms in the poles of ω



- * Computational bottleneck
- * Non-suitable for applications with finite-fields

● similar for multivariate case

$p(z)$ -adic expansion



$p(z)$ -adic series expansion

Expansion around all the roots of polynomials $p(z)$ at once

$$f(z) = \sum_{i=\min}^{\max} c_i(z) p^i(z) + \mathcal{O}(p(z)^{\max+1})$$

rational function

(prime) polynomial over \mathbb{Q}

polynomial coefficients $c_i(z)$

$$c_i(z) = \sum_{j=0}^{\deg p - 1} c_{ij} z^j$$

Obtained via repeated polynomial divisions



- * NO irrational operations
- * NO knowledge of explicit location

Example: univariate algorithm

Before:

$$\langle \varphi_L | \varphi_R \rangle = \sum_{p \in \mathcal{P}_\omega} \text{Res}_{z=p}(\psi \varphi_R)$$

summing over all $p \in \mathcal{P}_\omega$

$$\mathcal{P}_\omega = \{z \mid z \text{ is a pole of } \omega\} \cup \{\infty\}$$

Example: univariate algorithm

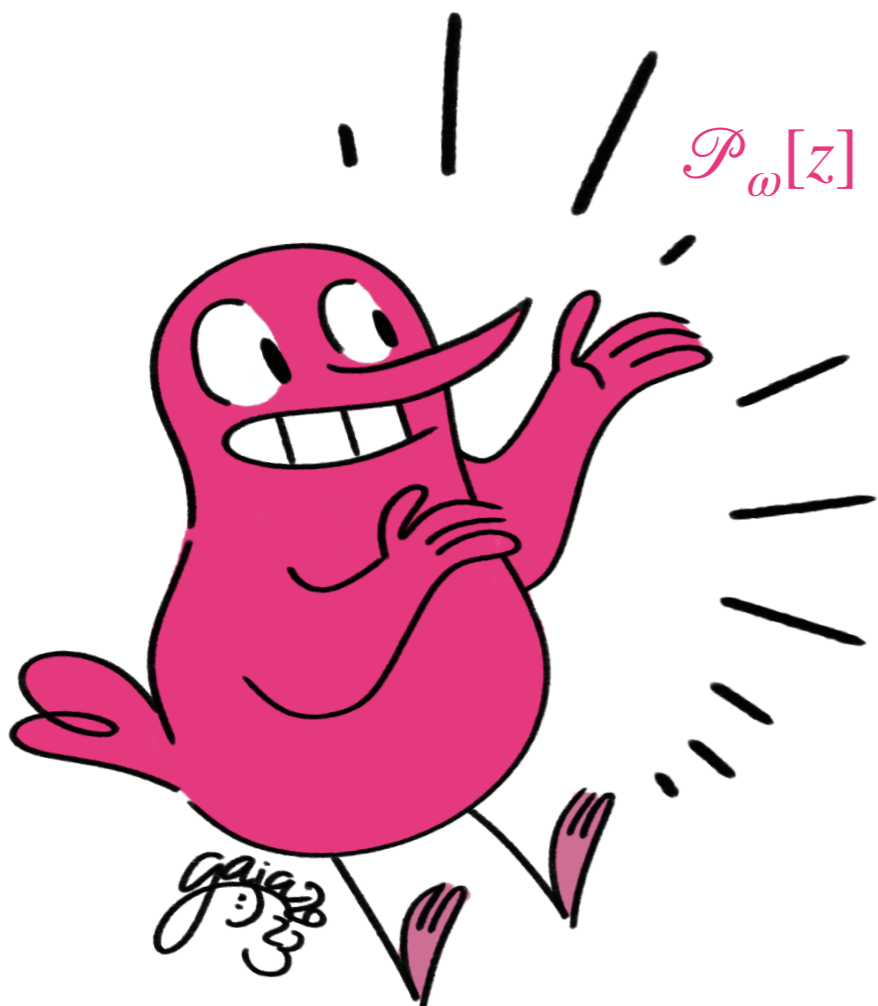
Now:

GF, T. Peraro (2022)

$$\langle \varphi_L | \varphi_R \rangle = \sum_{p(z) \in \mathcal{P}_\omega[z]} \langle \varphi_L | \varphi_R \rangle_{p(z)}$$

summing over all $p(z) \in \mathcal{P}_\omega[z]$

$\mathcal{P}_\omega[z] = \{\text{factors of the denominator of } \omega\} \cup \{\infty\}$



Example: univariate algorithm

GF, T. Peraro (2022)

$$\langle \varphi_L | \varphi_R \rangle = \sum_{p(z) \in \mathcal{P}_\omega[z]} \langle \varphi_L | \varphi_R \rangle_{p(z)}$$

summing over all $p(z) \in \mathcal{P}_\omega[z]$

$$\mathcal{P}_\omega[z] = \{\text{factors of the denominator of } \omega\} \cup \{\infty\}$$

- * Each addend of the form $\langle \varphi_L | \varphi_R \rangle_{p(z)}$ is the sum of all contributions to the intersection number coming from the roots of $p(z)$
- * $\langle \varphi_L | \varphi_R \rangle_\infty$ is computed as the contribution at $p = \infty$ with the “standard” algorithm

- similar for multivariate case

In practice:

to solve $(\partial_z + \omega)\psi = \varphi_L$

- we make an ansatz of the form

$$\psi = \sum_{i=\min}^{\max} \sum_{j=0}^{\deg p-1} c_{ij} z^j p(z)^i + O\left(p(z)^{\max+1}\right)$$

- we multiply the solution by φ_R

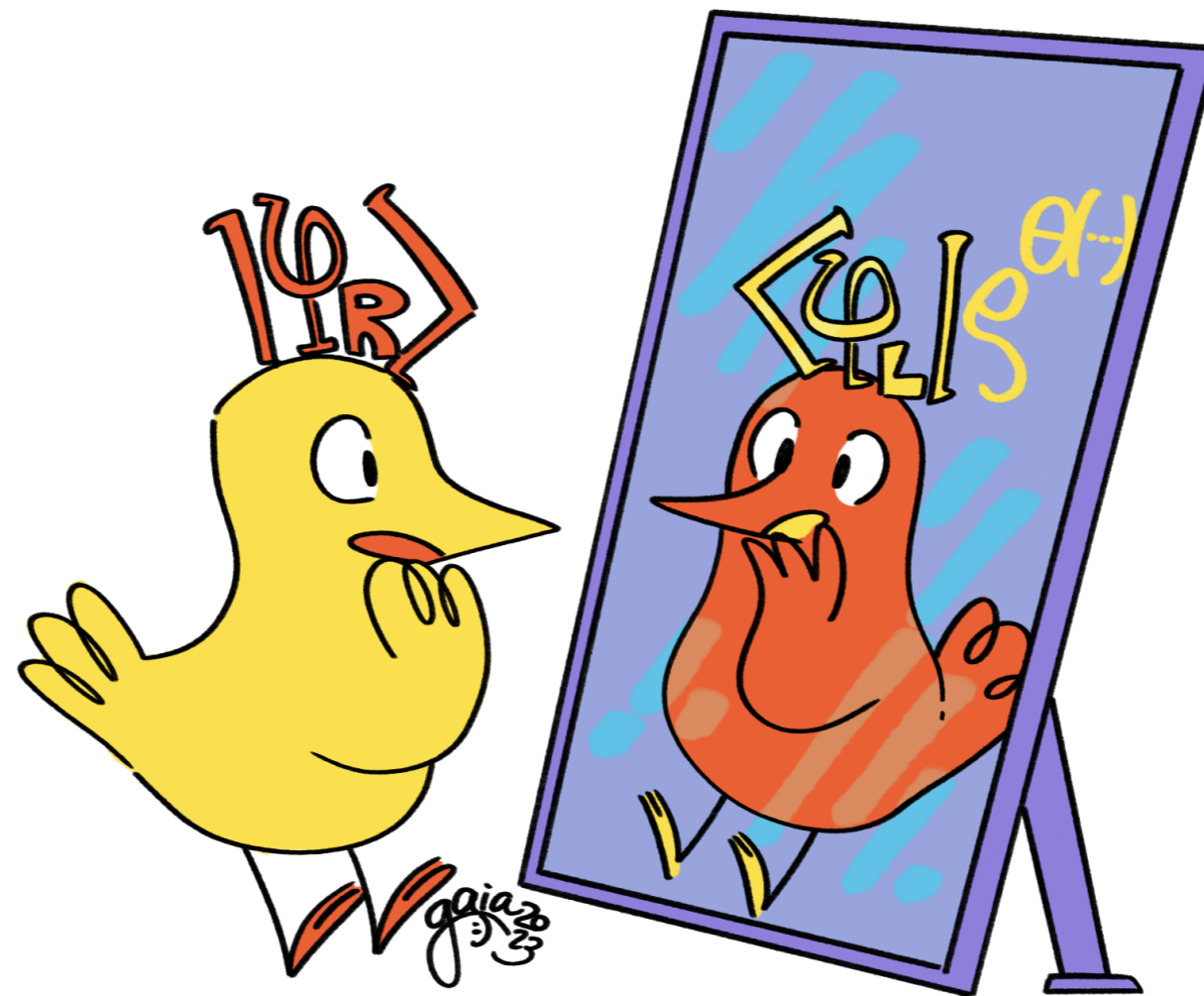
$$\psi\varphi_R = \sum_i^{-1} \sum_{j=0}^{\deg p-1} \tilde{c}_{ij} z^j p(z)^i + O\left(p(z)^0\right)$$

- by the univariate global residue theorem

Weinzierl (2021)

$$\langle \varphi_L | \varphi_R \rangle = \frac{\tilde{c}_{-1, \deg p-1}}{l_c}$$

Dual integrals



Analytic regulators: ρ_j

Are needed to regulate $\varphi_L \sim \frac{1}{z_j}$

$$u = B^{-\gamma} \prod_{j=1}^n z_j^{\rho_j} \quad \begin{array}{l} * \text{ otherwise no solution in DE for } \psi \\ * \text{ if } \varphi_L \sim 1/z^{v_j} \text{ (} v_j > 0 \text{) then } \psi \sim 1/\rho_j \end{array}$$

Then limit $\rho_j \rightarrow 0$ in the coefficients of the decomposition

BUT :(

- * Additional variables in intermediate steps
- * No block-triangular structure of decompositions
- * More MIs in intermediate steps

Dual Integrals...

Remember the right-integrals decomposition:

$$|\varphi_R\rangle = \sum_{i=1}^{\nu} c_i^{(R)} |e_i^{(R)}\rangle, \quad c_i^{(R)} = \sum_{j=1}^{\nu} (\mathbf{C}^{-1})_{ij} \langle e_j^{(L)} | \varphi_R \rangle$$

Observation

Coefficients $c_i^{(R)}$ are independent of the choice of the dual basis $\{\langle e_j^{(L)} | \}_{j=1}^{\nu}$

⇒ Idea

Exploit the freedom of choice of the dual basis to simplify the calculation

...& how to choose them

Two approaches (different formalism but similar outcomes)

- *Dual space of loop integrals [Caron-Huot, Pokraka (2021)]
- *Simple choice [GF, T. Peraro (2023)]

Choose dual integrals of the form


$$\varphi_L(\mathbf{z}) = \rho_1^{\Theta(\alpha_1 - \frac{1}{2})} \dots \rho_n^{\Theta(\alpha_n - \frac{1}{2})} \frac{1}{z_1^{\alpha_1} \dots z_n^{\alpha_n}}$$


- *If there's a denominator factor $z_j^{\alpha_j}$ (with $\alpha_j > 0$), multiply by ρ_j
- *Systematically work in the limit $\rho_j \rightarrow 0$ (only leading terms in a $\rho_j \rightarrow 0$ expansion in each step)




- * No dependence on ρ_j in calculations (work on leading coeff.s of $\rho_j \rightarrow 0$ expansion, never sample or reconstruct ρ_j dependence over FF)
- * **Simpler** intermediate expressions
- * **Block triangular** metric and reduction tables (blocks~sectors)

$$\mathbf{C} = \begin{pmatrix} * & * & 0 & 0 & 0 & 0 \\ * & * & 0 & 0 & 0 & 0 \\ * & * & * & * & 0 & 0 \\ * & * & * & * & 0 & 0 \\ * & * & 0 & 0 & * & * \\ * & * & 0 & 0 & * & * \end{pmatrix}$$

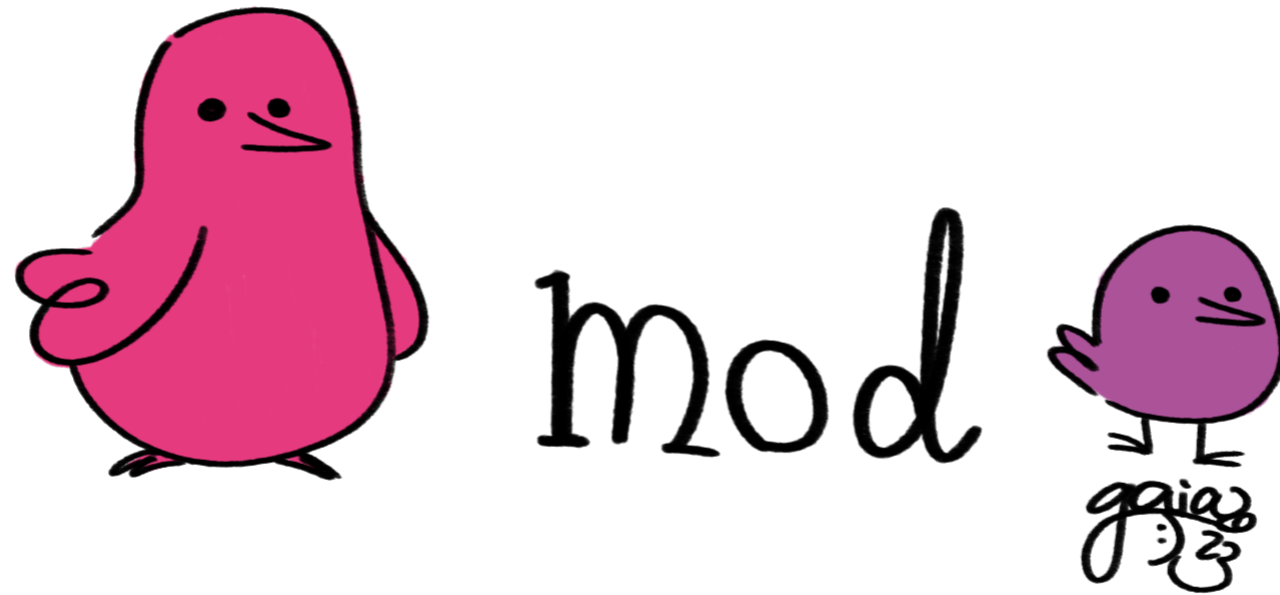
 top sector

 subsector 1

 subsector 2

- * Many intersection numbers and contributions of poles to them **vanish**
- * **Fewer MIs** in intermediate steps!

Finite fields implementation



Implementation on FiniteFlow of the multivariate recursive **rational** algorithm

GF, Peraro (2023)

● Input

list of n -variate intersection numbers to compute $\{\langle e_j^{(L)} | \varphi_R \rangle, \langle e_j^{(L)} | e_i^{(R)} \rangle\}$

● Preliminary step

recursively deduce the intersection numbers needed for each step

- * $\langle \varphi_L | e_j^{(R)} \rangle_{n-1}$
- * $(\partial_{z_n} \langle e_j^{(L)} |_{n-1} \rangle | e_j^{(R)} \rangle_{n-1}$
- * $\langle e_i^{(L)} | e_j^{(R)} \rangle_{n-1}$
- * $\langle e_j^{(L)} | \varphi_R \rangle$

● Univariate algorithm

analytic input: $u(\mathbf{z})$

● Multivariate algorithm

inputs

- * denominator factors $p_i(z_n)$
- * $(n-1)$ -variate intersection numbers reconstructed in z_n only

● Dealing with poles

- * $p = 0, \infty \rightarrow$ Laurent expansion
- * all other factors $\rightarrow p(z)$ -adic expansion

our implementation is an **iteration**:
1-forms $\rightarrow n$ -forms

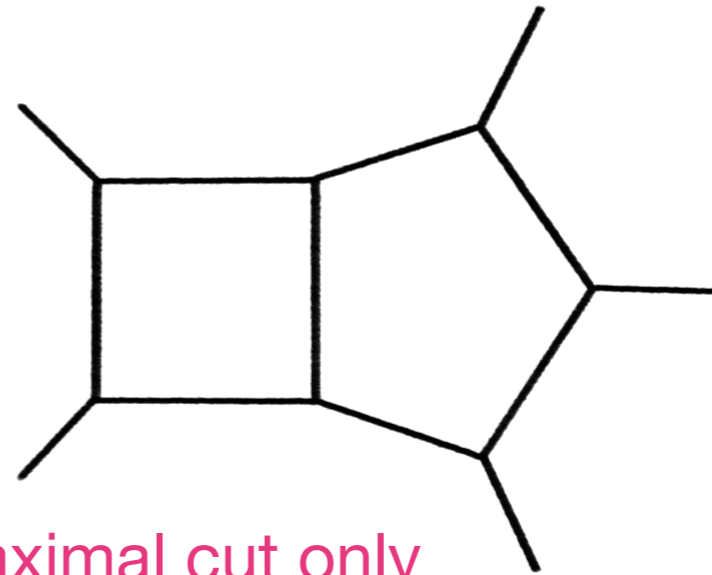
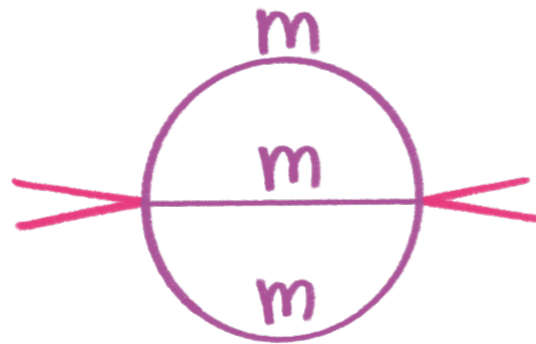
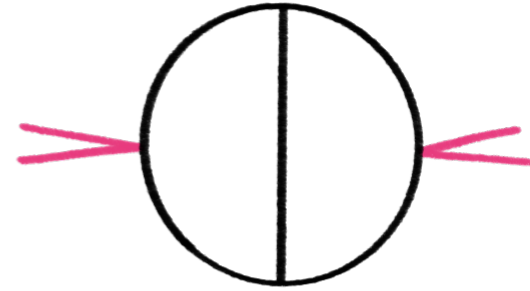
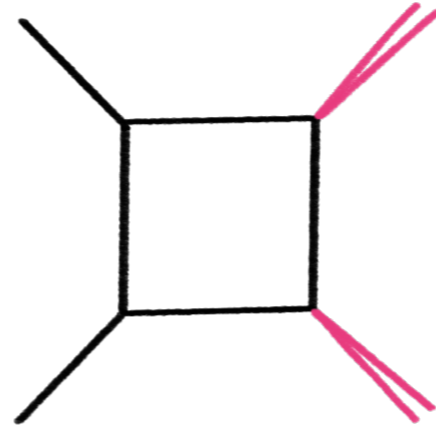
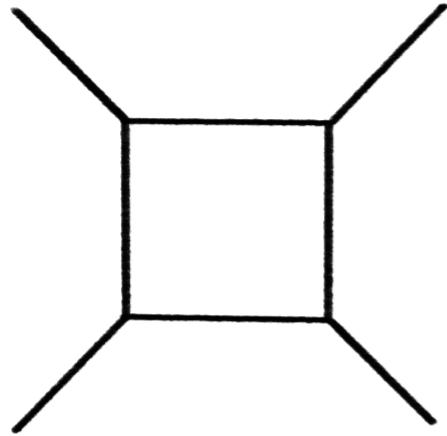
● Input for the n^{th} - step

$\mathcal{X}_n =$ list of $(n - 1)$ -variate intersection numbers and reduction coeff.s

● Between two steps:

- * rational reconstruction of \mathcal{X}_n only in z_n , with everything else set to a number mod p
- * identify denominator factors of \mathcal{X}_n in z_n , fully reconstruct them from a simple subset of \mathcal{X}_n

Examples



maximal cut only

Conclusions..

- * **Intersection theory:** new mathematical structures, direct integral reduction
- * **$p(z)$ -adic expansion:** simplify study of functions close to roots of polynomials

..& Outlook

- * Simplifications/optimizations
- * Application to different integral representations (loop-by-loop Baikov, Lee-Pomeransky)
- * Non-recursive multivariate generalization (based on [Chestnov, Frellesvig, Gasparotto, Mandal, Mastrolia \(2022\)](#))
- * New applications of $p(z)$ -adic expansion

Thank you for
your attention!

